

Lecture 25

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1 DFT, FFT, operations on polynomials

We are considering the complex polynomials of power $\leq n-1$. We will write this polynomial space as $V_{n-1}^{\mathbb{C}}$. We are interested to carry out on them 2 operations, multiplication, and value representation in time $O(n \log(n))$.

The two representations of a polynomial: coefficient representation, and value representation at n different complex points z_0, \dots, z_{n-1} .

The operation of value representation is efficient, and costs $O(n)$. In coefficient representation the multiplication operation is efficient in the value representation. We will focus on the value representation of the polynomials at the n roots of unity of power n .

Let us write $\omega_n = \cos\left(\frac{2\pi}{n}\right) + i \sin\left(\frac{2\pi}{n}\right)$. We will call ω_n the primitive root of unity of order n . So the powers ω_n^k where $0 \leq k \leq n-1$ are the roots of unity of order n . We will look at the values of the polynomial p at these points z_0, \dots, z_{n-1} where $z_k = \omega_n^k : 0 \leq k \leq n-1$.

Let there be $\begin{bmatrix} a_0 \\ \vdots \\ a_{n-1} \end{bmatrix}$ the coefficient vector of p , which is to say

$$p(z) = \sum_{k=0}^{n-1} a_k z^k$$

So $\begin{bmatrix} p(\omega_n^0) \\ \vdots \\ p(\omega_n^{n-1}) \end{bmatrix}$ will be called **the discrete fourier transform** of the vector $\begin{bmatrix} a_0 \\ \vdots \\ a_{n-1} \end{bmatrix}$, and we will write

$$\begin{bmatrix} p(\omega_n^0) \\ \vdots \\ p(\omega_n^{n-1}) \end{bmatrix} = DFT_n \begin{bmatrix} a_0 \\ \vdots \\ a_{n-1} \end{bmatrix}$$

Theorem 1 (FFT). *Let there be n , a power of two, and $\begin{bmatrix} a_0 \\ \vdots \\ a_{n-1} \end{bmatrix} \in \mathbb{C}^n$. Then we can find*

$$DFT_n \begin{bmatrix} a_0 \\ \vdots \\ a_{n-1} \end{bmatrix}$$

in $O(n \log(n))$.

Additionally, let there be $\begin{bmatrix} p_0 \\ \vdots \\ p_{n-1} \end{bmatrix} \in \mathbb{C}^n$. Then we can find $DFT_n^{-1} \begin{bmatrix} p_0 \\ \vdots \\ p_{n-1} \end{bmatrix}$ in $O(n \log(n))$

Proof intuition. We saw that the change between coefficient representation, and value representation of the polynomial p at the points z_0, \dots, z_{n-1} is the multiplying by the Vandermonde matrix. When the points z_0, \dots, z_{n-1} are the n th roots of unity, which is to say $0 \leq k \leq n-1$ $z_k = \omega_n^k$ we will get the matrix M (Vandermonde) such that

$$M = (z_k^m)_{k,m:0 \leq k,m \leq n-1} = (\omega_n^{km})_{k,m}$$

Example: $n = 4$: $\omega_n = i$, so $\omega_n^0 = 1$, $\omega_n^1 = i$, $\omega_n^2 = -1$, $\omega_n^3 = -i$. So

$$M_4 = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{bmatrix}$$

Lemma:

$$M_n^{-1} = \frac{1}{n} \cdot (\omega_n^{-km})_{k,m}$$

Also

$$\begin{bmatrix} A & B \\ A & -B \end{bmatrix} \begin{bmatrix} X \\ Y \end{bmatrix} = \begin{bmatrix} AX + BY \\ AX - BY \end{bmatrix}$$

$$T(n) \leq 2T\left(\frac{n}{2}\right) + O(n)$$

So if we count the rows and columns from 0, let us look for sub matrices of the shape

$$M_2 = \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

. So rows 0 and 2 of our matrix M_4 form 2 such matrices. Additionally from rows 1 and 3, we may take the matrices

$$B = \begin{bmatrix} 1 & i \\ 1 & -i \end{bmatrix}$$

$$-B = \begin{bmatrix} -1 & -i \\ -1 & i \end{bmatrix}$$

So we have split up M into the three required matrices, A which appears twice which is M_2 , and B and $-B$ □

Proof. Given a vector $\begin{bmatrix} a_0 \\ \vdots \\ a_{n-1} \end{bmatrix} \in \mathbb{C}^n$, (where n is a power of 2), we want to find $\begin{bmatrix} p(\omega_n^0) \\ \vdots \\ p(\omega_n^{n-1}) \end{bmatrix}$ in $O(n \log(n))$ (when

$$p(z) = \sum_{k=0}^{n-1} a_k z^k$$

Theorem 2 (Lemma). Let n be an even number, let $p(z) = \sum_{k=0}^{n-1} a_k z^k \in V_{n-1}$. We will define $p_0(y) = \sum_{j=0}^{\frac{n}{2}-1} a_{2j} y^j$ and

$$p_1(y) = \sum_{m=0}^{\frac{n}{2}-1} a_{2m+1} y^m. \text{ So}$$

$$p(z) = p_0(z^2) + z p_1(z^2)$$

Example: $n = 4$, $p(z) = 2z^3 - 3z^2 + z + 1$

$$\begin{aligned} 2z^3 - 3z^2 + z + 1 &= (-3z^2 + 1) + (2z^3 + z) \\ &= (-3z^2 + 1) + z \cdot (2z^2 + 1) \\ &= p_0(z^2) + z \cdot p_1(z^2) \\ p_0(y) &= -3y + 1 \\ p_1(y) &= 2y + 1 \end{aligned}$$

Proof .

$$\begin{aligned}
p_0(z^2) + zp_1(z^2) &= \sum_{j=0}^{\frac{n}{2}-1} a_{2j}(z^2)^j + z \cdot \sum_{m=0}^{\frac{n}{2}-1} a_{2m+1}(z^2)^m \\
&= \sum_{j=0}^{\frac{n}{2}-1} a_{2j}z^{2j} + \sum_{m=0}^{\frac{n}{2}-1} a_{2m+1}z^{2m+1} \\
&= \sum_{k=0}^{n-1} a_k z^k \\
&= p(z)
\end{aligned}$$

□

Theorem 3 (Lemma). *Let there be n an even number, then*

$$1. \omega_n^2 = \omega_{\frac{n}{2}}$$

2. For all $0 \leq j \leq \frac{n}{2} - 1$: $(\omega_n^j)^2 = \omega_{\frac{n}{2}}^j$, and also $\left(\omega_n^{\frac{n}{2}+j}\right)^2 = \omega_{\frac{n}{2}}^j$. In other words, when we raise the n th root of unity to the power of 2, we pass twice on the $\frac{n}{2}$ roots of unity

Proof . 1. $\omega_n = \cos\left(\frac{2\pi}{n}\right) + i \sin\left(\frac{2\pi}{n}\right)$. Therefore

$$\begin{aligned}
\omega_n^2 &= \cos\left(\frac{2\pi}{\frac{n}{2}}\right) + i \sin\left(\frac{2\pi}{\frac{n}{2}}\right) \\
&= \omega_{\frac{n}{2}}
\end{aligned}$$

2. For all $0 \leq j \leq \frac{n}{2} - 1$ it is true that

$$(\omega_n^j)^2 = \omega_n^{2j} = (\omega_n^2)^j = \omega_{\frac{n}{2}}^j$$

and so

$$\left(\omega_n^{\frac{n}{2}+j}\right)^2 = \omega_n^{2\left(\frac{n}{2}+j\right)} = \omega_n^n \cdot \omega_n^{2j} = \omega_{\frac{n}{2}}^j$$

□

We may now finally prove the theorem:

$$\begin{aligned}
\begin{bmatrix} p(\omega_n^0) \\ \vdots \\ p(\omega_n^{n-1}) \end{bmatrix} &= \begin{bmatrix} p_0((\omega_n^0)^2) + \omega_n^0 \cdot p_1((\omega_n^0)^2) \\ \vdots \\ p_0((\omega_n^{n-1})^2) + \omega_n^{n-1} \cdot p_1((\omega_n^{n-1})^2) \end{bmatrix} \\
&= \begin{bmatrix} p_0((\omega_n^0)^2) \\ \vdots \\ p_0((\omega_n^{n-1})^2) \end{bmatrix} + \begin{bmatrix} \omega_n^0 \\ \vdots \\ \omega_n^{n-1} \end{bmatrix} \bullet \begin{bmatrix} p_1((\omega_n^0)^2) \\ \vdots \\ p_1((\omega_n^{n-1})^2) \end{bmatrix} \quad \text{Multiplying coordinate - coordinate} \\
&= \begin{bmatrix} p_0\left(\omega_n^0\right) \\ \vdots \\ p_0\left(\omega_n^{\frac{n}{2}-1}\right) \\ p_0\left(\omega_n^{\frac{n}{2}}\right) \\ \vdots \\ p_0\left(\omega_n^{\frac{n}{2}-1}\right) \end{bmatrix} + \begin{bmatrix} \omega_n^0 \\ \vdots \\ \omega_n^{n-1} \end{bmatrix} \bullet \begin{bmatrix} p_1\left(\omega_n^0\right) \\ \vdots \\ p_1\left(\omega_n^{\frac{n}{2}-1}\right) \\ p_1\left(\omega_n^{\frac{n}{2}}\right) \\ \vdots \\ p_1\left(\omega_n^{\frac{n}{2}-1}\right) \end{bmatrix}
\end{aligned}$$

We will note that p_0 is a polynomial of power $\frac{n}{2} - 1$, with the coefficient vectors $\begin{bmatrix} a_0 \\ a_2 \\ \vdots \\ a_{n-2} \end{bmatrix}$, and therefore the vector

$$\begin{bmatrix} p_0\left(\omega_n^0\right) \\ \vdots \\ p_0\left(\omega_n^{\frac{n}{2}-1}\right) \end{bmatrix} \quad \text{which is the value vector of the polynomial } p_0, \text{ of the } \frac{n}{2} \text{ roots of unity.}$$

Therefore, from the definition of the DFT of order $\frac{n}{2}$

$$\begin{bmatrix} p_0\left(\omega_n^0\right) \\ \vdots \\ p_0\left(\omega_n^{\frac{n}{2}-1}\right) \end{bmatrix} = DFT_{\frac{n}{2}} \begin{bmatrix} a_0 \\ a_2 \\ \vdots \\ a_{n-2} \end{bmatrix}$$

and similarly

$$\begin{bmatrix} p_2\left(\omega_n^0\right) \\ \vdots \\ p_2\left(\omega_n^{\frac{n}{2}-1}\right) \end{bmatrix} = DFT_{\frac{n}{2}} \begin{bmatrix} a_1 \\ a_3 \\ \vdots \\ a_{n-1} \end{bmatrix}$$

and therefore we have the equivalence

$$DFT_n \begin{bmatrix} a_0 \\ \vdots \\ a_{n-1} \end{bmatrix} = \begin{bmatrix} DFT_{\frac{n}{2}} \begin{bmatrix} a_0 \\ a_2 \\ \vdots \\ a_{n-2} \end{bmatrix} \\ DFT_{\frac{n}{2}} \begin{bmatrix} a_0 \\ a_2 \\ \vdots \\ a_{n-2} \end{bmatrix} \end{bmatrix} + \begin{bmatrix} \omega_n^0 \\ \vdots \\ \omega_n^{n-1} \end{bmatrix} \bullet \begin{bmatrix} DFT_{\frac{n}{2}} \begin{bmatrix} a_1 \\ a_3 \\ \vdots \\ a_{n-1} \end{bmatrix} \\ DFT_{\frac{n}{2}} \begin{bmatrix} a_1 \\ a_3 \\ \vdots \\ a_{n-1} \end{bmatrix} \end{bmatrix}$$

□